Automatic generation of alternative starting positions for traditional board games

Umair Ahmed, Krishnendu Chatterjee, Sumit Gulwani

https://doi.org/10.15479/AT:IST-2013-146-v1-1

Deposited at: 12 Dec 2018 11:54
Automatic Generation of Alternative Starting Positions for Traditional Board Games

Umair Z. Ahmed and Krishnendu Chatterjee and Sumit Gulwani

Technical Report No. IST-2013-146-v1+1
Deposited at UNSPECIFIED
http://repository.ist.ac.at/146/1/main.pdf
Automatic Generation of Alternative Starting Positions for Traditional Board Games

Umair Z. Ahmed Krishnendu Chatterjee Sumit Gulwani

ABSTRACT

Board games, like Tic-Tac-Toe and CONNECT-4, play an important role not only in development of mathematical and logical skills, but also in emotional and social development. In this paper, we address the problem of generating targeted starting positions for such games. This can facilitate new approaches for bringing novice players to mastery, and also leads to discovery of interesting game variants. Our approach generates starting states of varying hardness levels for player 1 in a two-player board game, given rules of the board game, the desired number of steps required for player 1 to win, and the expertise levels of the two players. Our approach leverages symbolic methods and iterative simulation to efficiently search the extremely large state space. We present experimental results that include discovery of states of varying hardness levels for several simple grid-based board games. Also, the presence of such states for standard game variants like Tic-Tac-Toe on board size $4 \times 4$ opens up new games to be played that have not been played for ages since the default start state is heavily biased.

1. INTRODUCTION

Board games involve placing pieces on a pre-marked surface or board according to a set of rules by taking turns. Some of these grid-based two-player games like Tic-Tac-Toe and Connect-4 have a relatively simple set of rules, yet, they are recently challenging for certain age groups. Such games that are easy to learn but difficult to master have been immensely popular across centuries.

Significance of board games

Studies show that board games can significantly improve a child’s mathematical ability [16]. This is significant because studies show that differences in mathematical ability between children in the first year at school persist into secondary education [7]. Board games also play a vital role in the emotional and social development of a child. They instill a competitive urge and desire to master new skills in order to win. Winning gives a boost to their self-confidence and enjoyment. Playing a game within a set of rules helps them to adhere to discipline in life. They learn social etiquette; taking turns, and being patient. Strategy is another huge component of board games. Children should quickly grasp that decisions they make in the beginning of the game have consequences later on. Cause and effect is elegantly displayed in several board games.

Board games help elderly people stay mentally sharp and less likely to develop Alzheimer [11]. Board games also hold a great importance in today’s digital society by strengthening family ties. They bridge the gap between young and old. They bolster the self-esteem of children who take great pride and pleasure when an elder spends playing time with them.

Significance of generating fresh starting states

Board games are typically played with a default start state. For example, in case of Tic-Tac-Toe and CONNECT-4, it is the empty board. However, there are several drawbacks associated with starting from the default starting state. We highlight these problems below and accordingly motivate our goals.

Customizing hardness level of a start state. The default starting state for a certain game, while being unbiased, might not be conducive for a novice player to enjoy and learn the game. Traditional board games in particular are easy to learn but difficult to master because these games have intertwined mechanics and force the player to consider far too many possibilities from the standard starting configurations. Players can achieve mastery most effectively if complex mechanics can be simplified and learned in isolation.

Csikszentmihalyi’s theory of flow [6] suggests that we can keep the learner in a state of maximal engagement by continually increasing difficulty to match the learner’s increasing skill. Hence, we need an approach that allows generating start states of a specified hardness level. This capability can be used to generate a progression of starting states of increasing hardness. This is similar to how students are taught educational concepts such as addition and subtraction through a progression of increasingly hard problems [4].

Leveling the playing field. It is often the case that one player is more skilled than the other. The starting state for commonly played games is mostly unbiased, and hence does not offer a fair experience for players of different skills. The flexibility to start from some other starting state that is more biased towards the weaker player can allow for leveling the playing field and hence an enjoyable game among players with different expertise. Hence, we need an approach that takes as input the expertise levels of players and uses that information to associate a hardness level with a state.

Generating multiple fresh start states. A fixed starting state might have a well-known conclusion. For example, both players can enforce a draw in Tic-Tac-Toe while the first player can enforce a win in CONNECT-4 [3], starting from the default empty starting state. Players can memorize certain moves and strategies from a fixed starting state and gain undue advantage. Hence, we need an approach that generates multiple start states (of a specified hardness level). This observation has also inspired the design of Chess960 [1] (or Fischer Random Chess), which is a variant of chess that employs the same board and pieces as standard chess; however, the starting position of the pieces on the players’ home ranks is randomized. The random setup renders the prospect of obtaining an advantage through memorization of opening lines impracticable, compelling players to rely on their talent and creativity.

Customizing length of play. People sometimes might be disinterested in playing a game if it takes too much time to finish. However, selecting non-default starting positions allow the potential of a shorter game play. Certain interesting situations manifest only in
states that are typically not easily reachable from the start state, or might require too many steps. The flexibility to start from such states might lead to more opportunities for deliberate practice of specific targeted strategies. Hence, we need an approach that can take as input a parameter concerning the number of steps that can lead to a win for a given player.

Experimenting with game variants. While people might be hesitant to learn a game with completely different new rules, it is quite convenient to change the rules slightly. For example, instead of allowing for straight-line matches in each of row, column, or diagonal (RCD) in Tic-Tac-Toe or CONNECT-4, one may restrict the matches to say only row or diagonal (RD). However, the default starting state of a new game may be heavily biased towards a player; as a result that specific game might not have been popular. For example, consider the game of Tic-Tac-Toe (3,4,4), where the goal is to make a straight line of 3 pieces, but on a $4 \times 4$ board. In this game, the person who plays first invariably almost always wins even with a naive strategy. Hence, such a game has never been popular. It turns out that there are non-default unbiased states for such games and starting from those states can make playing such games interesting. Hence, we need an approach that is parameterized by the rules of a game. This also has the advantage of experimenting with new games or variants of existing games.

In summary, we need an approach to generate multiple starting states of specified hardness levels, given expertise levels of the players and length of plays, for traditional board games and their variants.

Problem Definition and Search Strategy

We address the problem of automatically generating interesting starting states (i.e., states of desired hardness levels) for a given two-player board game. Our approach takes as input the rules of a board game (for game variants) and the desired number of steps required for player 1 to win (for controlling the length of play). It then generates multiple starting states of varying hardness levels (in particular, easy, medium, or hard) for player 1 for various expertise level combinations of the two players.

In our setting of simple board games, the players fully know the simple rules and all the legal game states (as opposed to the setting of complicated games, where players construct, with approximation and error, rules of the game and search in their own hypothesized state space). Hence, we formalize the exploration of a game as a strategy tree and the expertise level of a player as depth of the strategy tree. The hardness of a state is defined with respect to the fraction of times player 1 will win, while playing a strategy of depth $k_1$ against an opponent who plays a strategy of depth $k_2$.

Our solution is a novel combination of symbolic methods and iterative simulation to efficiently search for desired states. Symbolic methods are used to compute the winning set for player 1. These methods work particularly well for navigating a state space where the transition relation forms a sparse directed acyclic graph (DAG). Such is the case for those board games in which a piece once placed on the board does not move, as in Tic-Tac-Toe and CONNECT-4. Minimax simulation is used to identify the hardness of a given winning state. Instead of randomly sampling the winning set to identify a state of a certain hardness level, we identify states of varying hardness levels in order of increasing values of $k_1$ and $k_2$. The key observation is that hard states are much smaller in number than easy states, and given a value of $k_2$, interesting states for higher values of $k_1$ are a subset of the hard states for smaller values of $k_1$.

Results

Though our general search methodology applies for any graph game, we focus on generating interesting starting states in traditional simple board games and their variants—these are games whose transition relation forms a sparse DAG (as opposed to an arbitrary graph). Generating starting states in simple and traditional games, as compared to games with complicated rules, is both more challenging and more relevant. First, in sophisticated games with complicated rules interesting states are likely to be abundant and hence easier to find, whereas finding interesting states in simple games is more challenging. Second, games with complicated rules are hard to learn, whereas simple variants of traditional games (such as larger or smaller board size, or changing winning conditions from RCD to RD) are easier to adopt. Hence finding interesting start states in traditional games and its variants is the more relevant and challenging question that we address in this work. We thus present a framework to easily describe new board games like Tic-Tac-Toe or CONNECT-4 or their variants. Our implementation works for games that can be described using this framework.

We experimented with Tic-Tac-Toe, CONNECT, Bottom-2 (a new game that is a hybrid of Tic-Tac-Toe and CONNECT) games and several of their variants like RD or RC as winning rules instead of RCD. We were able to generate several (ten of) starting states of various hardness levels for various expertise levels and number of winning steps. Two important findings of the experiments are: (i) discovery of starting states of various hardness levels in these traditional board games, and furthermore discovering them in games such as Tic-Tac-Toe on $4 \times 4$ board size where the default start state is heavily biased; and (ii) these states are rare and thus require a non-trivial search strategy like ours to find them.

Contributions

This paper makes the following contributions.

- We motivate and formalize the problem of generating starting states for board games, and in general, graph games. In particular, we formalize the problem of generating starting states of varying hardness levels parameterized by the expertise levels of players, the game graph description, and the number of steps required for winning (§2).
- We present a novel search methodology for generating desired initial states. It involves combination of symbolic methods and iterative simulation to efficiently search a huge state space (§3).
- We present experimental results that illustrate the effectiveness of our search methodology (§5). We produced a collection of initial states of varying hardness levels for standard games as well as their variants (thereby discovering some interesting variants of the standard games in the first place).

2. Problem Definition

In this section, we first present some necessary background related to mathematical model of graph games and recall some basic results (§2.1). We then describe the notion of hardness and finally the description of the problem we consider for graph games (§2.2).

2.1 Background on Graph Games

Graph games. An alternating graph game (for short, graph game) $G = ((V,E), (V_1, V_2))$ consists of a finite graph $G$ with vertex set $V$, a partition of the vertex set into player-1 vertices $V_1$ and player-2 vertices $V_2$, and edge set $E \subseteq ((V_1 \times V_2) \cup (V_2 \times V_1))$. The game is alternating in the sense that the edges of player-1 vertices go to player-2 vertices and vice-versa. The game is played as follows: the game starts at a starting vertex $v_0$; if the current vertex is a player-1 vertex, then player 1 chooses an outgoing edge to move to a new vertex; if the current vertex is a player-2 vertex, then player 2 does
likewise. The winning condition is given by a target set $T_1 \subseteq V$ for player 1; and similarly a target set $T_2 \subseteq V$ for player 2. If the target set $T_1$ is reached, then player 1 wins; if $T_2$ is reached, then player 2 wins; else we have a draw.

**Examples.** The class of graph games provides the mathematical framework to study many board games like Chess or Tic-Tac-Toe. For example, in Tic-Tac-Toe the vertices of the graph represent the board configurations and whether it is player 1's (×) or player 2's (○) turn to play next. The set $T_1$ (resp. $T_2$) is the set of board configurations with three consecutive × (resp. ○) in a row, column, or diagonal.

**Classical game theory result.** A classic result in the theory of graph games [8] shows that for every graph game, from every starting vertex one of the following three conditions hold: (1) player 1 can enforce a win no matter how player 2 plays (i.e., there is a way for player 1 to play to ensure winning against all possible strategies of the opponent); (2) player 2 can enforce a win no matter how player 1 plays; or (3) both players can enforce a draw (player 1 can enforce a draw no matter how player 2 plays, and player 2 can enforce a draw no matter how player 1 plays). The classic result (also known as determinacy) rules out the following possibility: against every player 1 strategy (way to play), player 2 can win; and against every player 2 strategy, player 1 can win. In the mathematical study of game theory, the theoretical question (which ignores the notion of hardness) is as follows: given a designated starting vertex $v_0$ determine whether case (1), case (2), or case (3) holds. In other words, the mathematical game theoretic question concerns the best possible way for a player to play to ensure the best possible result. The set $W_1$ is defined as the set of vertices such that player 1 can ensure to win within $j$-moves; and the winning set $W^1$ of vertices of player 1 is the set $\bigcup_{j \geq 0} W_j$ where player 1 can win in any number of moves. Analogously, we define $W_2^2$; and then the classical game theory question is formally stated as follows: given a designated starting vertex $v_0$ decide whether $v_0$ belongs to $W^1$ (player-1 winning set) or to $W^2$ (player-2 winning set) or to $V \setminus (W^1 \cup W^2)$ (both players draw ensuring set).

**2.2 Formalization of Problem Definition**

We state some useful terminology and then our problem definition.

**Notion of hardness.** The mathematical game theoretic question ignores two aspects. (1) The notion of hardness: It is always concerned with optimal strategies irrespective of its hardness, and it is not concerned with when can sub-optimal strategies perform well, and when do sub-optimal strategies fail; and (2) the problem of generating different starting vertices. In this work we are interested in generating starting vertices of different hardness level. The notion of hardness we consider is the depth of the tree a player can explore, which is standard in artificial intelligence. We first explain the notion of tree exploration.

**Tree exploration in graph games.** Consider a player-1 vertex $u_0$. The search tree of depth 1 is as follows: we consider a tree rooted at $u_0$ such that children of $u_0$ are the vertices $v_1$ of player 2 such that $(u_0, v_1) \in E$ (there is an edge from $u_0$ to $u_1$); and for every vertex $u_2$ (that is a children of $u_0$) the children of $u_1$ are the vertices $v_2$ such that $(u_1, v_2) \in E$; and they are the leaves of the tree. This gives us the search tree of depth 1, which intuitively corresponds to exploring one round of the play. The search tree of depth $k + 1$ is defined inductively from the search tree of depth $k$, where we first consider the search tree of depth 1 and replace every leaf by a search tree of depth $k$. The depth of the search tree denotes the depth of reasoning (analysis depth) of a player and corresponds to the optimality (or competence or maturity) of the player to play the game. The search tree for player 2 is defined analogously.

**Strategy from tree exploration.** A depth-$k$ strategy of a player that does a tree exploration of depth $k$ is obtained by the classical min-max reasoning (or backward induction) on the search tree. First, for every vertex $v$ of the game we associate a number (or reward) $r(v)$ that denotes how favorable is the vertex for a player to win. Given the current vertex $u$, a depth-$k$ strategy is defined as follows: first construct the search tree of depth $k$ and evaluate the tree bottom-up with min-max reasoning. In other words, a leaf vertex is assigned reward $r(v)$, where the reward function $r$ is game specific, and intuitively, $r(v)$ denotes how “close” the vertex $v$ is to a winning vertex (see the following paragraph for an example). For a vertex in the tree if it is a player-1 (resp. player-2) vertex we consider its reward as the maximum (resp. minimum) of its children, and finally, for vertex $v$ (the root) the strategy chooses uniformly at random among its children with the highest reward. Note that the rewards are assigned to vertices only based on the vertex itself and without any look-ahead, and the exploration (or look-ahead) is captured by the classical min-max tree exploration.

**Example description of tree exploration.** Consider the example of the Tic-Tac-Toe game. We first describe how to assign reward $r$ to board positions. Recall that in the game of Tic-Tac-Toe the goal is to form a line of three consecutive positions in a row, column, or diagonal. Given a board position, (i) if it is winning for player 1, then it is assigned reward $+\infty$; (ii) else if it is winning for player 2, then it is assigned reward $-\infty$; (iii) otherwise it is assigned the score as follows: let $n_1$ (resp. $n_2$) be the number of two consecutive positions of marks for player 1 (resp. player 2) that can be padded to satisfy the winning condition. The difference $n_1 - n_2$ is the difference of the number of possibilities for player 1 to win, and $n_2$ represents the number of possibilities for player 2, and their difference represents how favorable the board position is for player 1. If we consider the depth-1 strategy, then the strategy chooses all board positions uniformly at random; a depth-2 strategy chooses the center and considers all other positions to be equal; a depth-3 strategy chooses the center and also recognizes that the next best choice is one of the four corners. An illustration is given in the appendix. This example illustrates that as the depth increases, the strategies become more intelligent for the game.

**Outcomes and probabilities given strategies.** Given a starting vertex $v$, a depth-$k$, strategy $\sigma_v$ for player 1, and depth-$k_2$ strategy $\sigma_{v_2}$ for player 2, let $O$ be the set of all possible outcomes, i.e., the set of all possible plays given $\sigma_1$ and $\sigma_2$ from $v$, where a play is a sequence of vertices. The strategies and the starting vertex define a probability distribution over the set of outcomes which we denote as $Pr^{\sigma_1, \sigma_2}(\rho)$, i.e., for a play $\rho$ in the set of outcomes $O$ we have $Pr^{\sigma_1, \sigma_2}(\rho)$ is the probability of $\rho$ given the strategies. Note that strategies are randomized (because strategies choose distributions over the children in the search tree exploration), and hence define a probability distribution over the set of outcomes. The notion of probability distribution will be used to formally define the notion of hardness we consider.

**Problem definition.** In this work we will consider several board games (such as Tic-Tac-Toe, CONNECT-4, and variants), and the goal is to obtain starting positions that are of different hardness level, where our hardness is characterized by strategies of different depths. In other words, our goal is to obtain starting positions that are of different hardness levels: i.e., hard for depth-1 strategies, but easy for depth-2 strategies; and hard for depth-2 strategies, but easy for depth-3 strategies, and so on. More precisely, consider a depth-$k_1$ strategy for player 1, and depth-$k_2$ strategy for player 2, and a starting vertex $v \in W_1$ that is winning for player 1 within $j$-moves and a winning move (i.e., $j + 1$ moves for player 1 and $j$ moves of
player 2). We classify the starting vertex into three types as follows: if player 1 wins (i) at least \( \frac{1}{2} \) times, then we call it easy (\( E \)); (ii) at most \( \frac{1}{2} \) times, then we call it hard (\( H \)); (iii) otherwise medium (\( M \)).

**Definition 1.** \((k_i, k_j, k_2)\)-Hardness. Consider a vertex \( v \in \mathcal{V}_j \) that is winning for player 1 within \( j \)-moves. Let \( \sigma_1 \) and \( \sigma_2 \) be a depth-\( k_1 \) strategy for player 1 and depth-\( k_2 \) strategy for player 2, respectively. Let \( O_1 \subseteq \mathcal{O}_1 \) be the set of plays that belong to the set of outcomes and is winning for player 1. Let \( \Pr^{v_1 \rightarrow \sigma_2}(O_1) = \sum_{\rho \in O_1} \Pr^{v_1 \rightarrow \sigma_2}(\rho) \) be the probability of the winning plays. The \((k_1, k_2)\)-classification of \( v \) is as follows: (i) if \( \Pr^{v_1 \rightarrow \sigma_2}(O_1) \geq \frac{3}{4} \), then \( v \) is easy (\( E \)); (ii) if \( \Pr^{v_1 \rightarrow \sigma_2}(O_1) \leq \frac{1}{4} \), then \( v \) is hard (\( H \)); (iii) otherwise it is medium (\( M \)).

**Remark.** In the definition above we chose the probabilities \( \frac{1}{4} \) and \( \frac{3}{4} \), however, the probabilities in the definition could be easily changed and experimented. We chose \( \frac{1}{4} \) and \( \frac{3}{4} \) to divide the interval \([0, 1]\) symmetrically in regions of \( E, M, \) and \( H \). In this work, we present result based on the above definition.

Our goal is to consider various games and identify states of different categories (e.g., hard for depth-\( k_1 \) against depth-\( k_2 \), but easy for depth-(\( k_1+1 \)) against depth-\( k_2 \), for small values of \( k_1 \) and \( k_2 \)).

### 3. SEARCH STRATEGY

We now describe the search strategy that we use for the problem of generating starting positions of different hardness levels.

#### 3.1 Overall methodology

We start with a description of the overall methodology.

**Generation of \( j \)-steps win set.** Given a game graph \( G = (\mathcal{V}, \mathcal{E}, (\mathcal{V}_1, \mathcal{V}_2)) \) along with target sets \( \mathcal{T}_1 \) and \( \mathcal{T}_2 \) for player 1 and player 2, respectively, our first goal is to compute the set of vertices \( \mathcal{W}_j \) such that player 1 can win within \( j \)-moves. For this we define two kinds of predecessor operators: one predecessor operator for player 1, which uses existential quantification over successors, and one for player 2, which uses universal quantification over successors. Given a set of vertices \( \mathcal{X} \), let \( \mathcal{EPre}(\mathcal{X}) \) (called existential predecessor) denote the set of player 1 vertices that has an edge to \( \mathcal{X} \); i.e., \( \mathcal{EPre}(\mathcal{X}) = \{ u \in \mathcal{V}_1 \mid \text{there exists} v \in \mathcal{X} \text{ such that} (u, v) \in \mathcal{E} \} \) (i.e., player 1 can ensure to reach \( \mathcal{X} \) from \( \mathcal{EPre}(\mathcal{X}) \) in one step); and \( \mathcal{APre}(\mathcal{X}) \) (called universal predecessor) denote the set of player 2 vertices that has all its outgoing edges to \( \mathcal{X} \); i.e., \( \mathcal{APre}(\mathcal{X}) = \{ u \in \mathcal{V}_2 \mid \text{for all} (u, v) \in \mathcal{E} \text{ we have} v \in \mathcal{X} \} \) (i.e., irrespective of the choice of player 2 the set \( \mathcal{X} \) is reached from \( \mathcal{APre}(\mathcal{X}) \) in one step). The computation of the set \( \mathcal{W}_j \) is defined inductively as follows: \( \mathcal{W}_0 = \mathcal{EPre}(\mathcal{T}_1) \) (i.e., player 1 wins with the next move to reach \( \mathcal{T}_1 \)); and \( \mathcal{W}_{j+1} = \mathcal{APre}(\mathcal{EPre}(\mathcal{W}_j)) \). In other words, from \( \mathcal{W}_j \), player 1 can win within \( j \)-moves, and from \( \mathcal{APre}(\mathcal{W}_j) \) irrespective of the choice of player 2 the next vertex is in \( \mathcal{W}_j \); hence \( \mathcal{EPre}(\mathcal{APre}(\mathcal{W}_j)) \) is the set of vertices such that player 1 can win within \((j+1)\)-moves.

**Exploring vertices from \( \mathcal{W}_j \).** The second step is to explore vertices from \( \mathcal{W}_j \), for increasing values of \( j \) starting with small values of \( j \), for strategies of different depth for player 1 against strategies of different depth for player 2, and then obtain starting vertices of different hardness levels. That is, we consider a vertex \( v \in \mathcal{W}_j \), consider a depth-\( k_1 \) strategy for player 1 and a depth-\( k_2 \) strategy for player 2, and play the game multiple times starting vertex \( v \) to find out the hardness level with respect to \((k_1, k_2)\)-strategies, i.e., the \((k_1, k_2)\)-classification of the vertex \( v \in \mathcal{W}_j \). Note that from \( \mathcal{W}_j \), player 1 can win within \( j \)-moves. Thus the approach has the benefit that player 1 has a winning strategy with a small number of moves and the game need not be played for long.

**Two key issues.** There are two main computational issues associated with the above approach in practice. The first issue is related to the size of the state space (number of vertices) of the game which makes enumerative approach to analyze the game graph explicitly computationally infeasible. For example, the size of the state space of a Tic-Tac-Toe \( 3 \times 3 \) game is 5,478; Tic-Tac-Toe \( 4 \times 4 \) game is 6,036,001; and a CONNECT-\( 5 \times 5 \) game is 69,763,700 (more than 69 million). Hence any enumerative method would not work for such large game graphs. The second issue is related to exploring the vertices from \( \mathcal{W}_j \). If \( \mathcal{W}_j \) has a lot of witness vertices, then playing the game multiple times from all of them will be computationally expensive. In other words, we need an initial metric to guide the search of vertices from \( \mathcal{W}_j \) such that the metric computation is not computationally expensive. We solve the first issue with symbolic methods, and the second one by iterative simulation.

#### 3.2 Symbolic methods

In this section we discuss the symbolic methods that allow to analyze games with large state spaces. The key idea is to represent the games symbolically (not with explicit state space) using variables, and operate on the symbolic representation. The key object used in symbolic representation are called BDDs (boolean decision diagrams) [5] that can efficiently represent a set of states using a dag representation of a boolean formula representing the set of states. The tool CuDD supports many symbolic representation of state space using BDDs and supports many operations on symbolic representation on graphs using BDDs [20].

**Symbolic representation of vertices.** In symbolic methods, a game graph is represented by a set of variables \( x_1, x_2, \ldots, x_n \), such that each of them takes values from a finite set (e.g., \( \{x, o, \text{blank}\} \)) and each vertex of the game represents a valuation assigned to the variables. For example, the symbolic representation of the game of Tic-Tac-Toe of board size \( 3 \times 3 \) consists of ten variables \( x_{1,1}, x_{1,2}, x_{1,3}, x_{2,1}, \ldots, x_{3,3}, x_{3,9} \), where the first nine variables \( x_{i,\ell} \) denote the symbols in the board position \((i, \ell)\) and the symbol is either \( x, o, \) or blank; and the last variable \( x_{3,9} \) denote whether it is player 1 or player 2’s turn to play. Note that the vertices of the game graph does not only consists of the information about the board configuration, but also additional information such as the turn of the players. To illustrate how a symbolic representation is efficient, consider the set of all valuations to boolean variables \( y_1, y_2, \ldots, y_n \) where the first variable is true, and the second variable is false: an explicit enumeration requires to list \( 2^{n-2} \) valuations, where as a boolean formula representation is very succinct. Symbolic representation with BDDs exploit such succinct representation for sets of states, and widely used in many applications such as hardware verification [5].

**Symbolic encoding of transition function.** The transition function (or the edges) are also encoded in a symbolic fashion: instead of specifying every edge, the symbolic encoding allows to write a simple program over the variables to specify the transitions. The tool CuDD takes such a symbolic description written as a program over the variables and constructs a BDD representation of the transition function. For example, for the game of Tic-Tac-Toe, a simple program to describe the symbolic transition is as follows: the program maintains a set \( \mathcal{U} \) of positions of the board that are already marked; and at every point receives an input \((i, \ell)\) from the set \( \{(a, b) \mid 1 \leq a, b \leq 3\} \) \( \setminus \mathcal{U} \) of remaining board positions from the player of the current turn; then adds \((i, \ell)\) to the set \( \mathcal{U} \) and sets the variable \( x_{i,\ell} = x \) or \( o \) (depending on whether it was player 1 or player 2). This represents the symbolic description of the transition function. The CuDD tool accepts such a symbolic description and outputs a BDD representation of the game.
Symbolic encoding of target states. The set of target states is encoded as a boolean formula that represents a set of states. For example, in Tic-Tac-Toe the set of target vertices for player 1 is given by the following boolean formula:

\[
\exists i,l. 1 \leq i,l \leq 3, (x_{i,l} = x \land x_{i+1,l} = x \land x_{i+2,l} = x) \land \neg (x_{i,l} = x \land x_{i,l+1} = x) \land \neg (x_{i,l} = x \land x_{i+2,l} = x)
\]

The above formula states that either there is some column \((x_{i,l}, x_{i+1,l}, x_{i+2,l})\) that is winning for player 1; or a row \((x_{i,l}, x_{i,l+1}, x_{i,l+2})\) that is winning for player 1; or there is a diagonal \((x_{1,1}, x_{2,2}, x_{3,3}), (x_{1,3}, x_{2,2}, x_{3,1}), (x_{1,1}, x_{2,2}, x_{3,3})\) that is winning for player 1; and player 2 has not won already. To be precise, we also need to consider the BDD that represents all valid board configurations (reachable vertices from the empty board) and intersect the BDD of the above formula with valid board configurations to obtain the target set \(T_1\).

Symbolic computation of \(W_j\). The symbolic computation of \(W_j\) is as follows: given the boolean formula for the target set \(T_1\) we obtain the BDD for \(T_1\); and the CuDD tool supports both EPre and APRe as basic operations using symbolic functions; i.e., the tool takes as input a BDD representing a set \(X\) and supports the operation to return the BDD for \(EPre(X)\) and \(APRe(X)\). Thus we obtain the symbolic computation of the set \(W_j\).

### 3.3 Iterative simulation

We now describe a computationally inexpensive way to aid sampling of vertices as candidates for starting positions of a given hardness level. Given a starting vertex \(v\), a depth-\(k_1\) strategy for player 1, and a depth-\(k_2\) strategy for player 2, we need to consider the tree exploration of depth \(\max(k_1, k_2)\) to obtain the hardness of \(v\). Hence if either of the strategy is of high depth, then it is computationally expensive. Thus we need a preliminary metric that can be computed relatively easily for small values of \(k_1\) and \(k_2\) as a guide for vertices to be explored in depth. We use a very simple metric for this purpose. The hard states are quite rare in comparison to the easy states, and thus we need to rule out easy states quickly. We adopt the following two approaches.

- **If \(k_1\) is large.** Given a strategy of depth \(k_2\), the set of hard states for higher values of \(k_1\) are a subset of the hard states for smaller values of \(k_1\). Thus we iteratively start with smaller values of \(k_2\) and proceed to higher values of \(k_1\) only for states that are hard already for smaller values of \(k_1\).

- **If \(k_2\) is large.** Here we exploit the following intuition. Given a strategy of depth \(k_1\), a state which is hard for high value of \(k_2\) is likely to show indication of hardness already in small values of \(k_2\). Hence we consider the following approach. For the vertices in \(W_j\), we fix a depth-\(k_1\) strategy, and fix a small depth strategy for the opponent and assign the vertex a number (called score) based on the performance of the depth-\(k_1\) strategy and the small depth strategy of the opponent. The score indicates the fraction of games won by the depth-\(k_1\) strategy against the opponent strategy of small depth. The vertices that have low score are then iteratively simulated against depth-\(k_2\) strategies of the opponent to obtain vertices of different hardness level. This heuristic serves as a simple metric to explore vertices for large value of \(k_2\) starting with small values of \(k_2\).

### 4. FRAMEWORK FOR BOARD GAMES

We now consider the specific problem of board games. We describe a framework to specify several variants of two-player grid based board games such as Tic-Tac-Toe, CONNECT-4, and several new variants. Note that though our implementation of symbolic methods works for the class of traditional board games and their variants, our methodology is applicable to the general class of graph games.

#### Parameters to generate different games.

Our framework allows three different parameters to generate variants of board games.

1. **The first parameter is the board size:** e.g., the board size could be \(3 \times 3\); or \(4 \times 4\); or \(4 \times 5\) and so on.

2. **The second parameter is the operation to specify the winning condition:** and we consider the cases where a player wins if a sequence of the moves of the player are in a line but the line could be in a row (R), in a column (C), or along the diagonal (D). The user can specify any combination, i.e., the winning condition could be (i) RCD (denoting the player wins if the moves are in a line along a row, column, or diagonal); (ii) RC (line must be along a row or column, but diagonal lines are not winning); (iii) RD (row or diagonal, but not column); or (iv) CD (column or diagonal, but not row).

3. **The third parameter is related to the allowed moves of the player.** At any point the players can choose a column (if it is available, i.e., there is at least one empty position in the column) but can be restricted according to the following parameters: (i) Full gravity (once a player chooses a column, the move is fixed to be the lowest available position in that column); (ii) partial gravity-\(\ell\) (once a player chooses a column, the move can be one of the \(\ell\) available positions in the column); (iii) no gravity (the player can choose any of the available positions in the column).

Observe that Tic-Tac-Toe is given as board size \((i)\) board size \(3 \times 3\); (ii) winning condition RCD; and (iii) no-gravity; whereas in CONNECT-4 the winning condition is still RCD but moves are with full gravity. But in our framework there are many new variants of the previous classical games, e.g., Tic-Tac-Toe in a board of size \(4 \times 4\) but diagonal lines are not winning (RC). Tic-Tac-Toe, Bottom-2 and CONNECT-3 require 3 consecutive positions to be marked for a player to win, while CONNECT-4 requires 4 consecutive positions. Note that the Bottom-2 (partial gravity-2) is between Tic-Tac-Toe and CONNECT games in terms of moves allowed: in Tic-Tac-Toe all available positions are allowed, whereas in CONNECT games a player can choose among the available columns, and in Bottom-2 a player can choose an available column and among the two bottom positions available in the column.

#### Features of our implementation.

We have implemented our symbolic approach for generating starting vertices (or board positions) of different hardness levels (if they exist) for the class of board games described above. The main features that our implementation supports are: (1) Generation of starting vertices of different hardness level if they exist. (2) Playing against opponents of different levels. We have implemented the depth-\(k_2\) strategy of the opponent for \(k_2 = 1, 2, 3\) (typically in all the above games depth-3 strategies are quite intelligent, and hence we do not explore larger values of \(k_2\)). Thus, a learner (beginner) can consider starting with board positions of various hardness levels and play with opponents of different skill level and thus hone her ability to play the game and be exposed to new combinatorial challenges of the game.

#### 5. EXPERIMENTAL RESULTS

We now present our experimental results, which reveal useful discoveries. The main aim is to investigate the existence of interesting starting states and their abundance (if they exist) in CONNECT, Tic-Tac-Toe, and Bottom-2 games, for various combinations of expertise levels and for various winning rules (RCD, RC, RD, and CD), for small lengths of plays. Moreover, the experimental results should be accomplished in reasonable time. Our key findings
Figure 1: Some “Hard” starting board positions generated by our tool for a variety of games and a variety of expertise level \(k_1\) of the first player. The opponent expertise level \(k_2\) is set to 3. The first player (player X) can win in 2 steps for games (a)-(e) and in 3 steps for game (f).

Table 1: CONNECT-3 & -4 against depth-3 strategy of opponent; (C-3 (resp. C-4) stands for CONNECT-3 (resp. CONNECT-4)). The third column \(j\) denotes whether we explore from \(W_2\) or \(W_3\). The sixth column denotes sampling to select starting vertices if \(|V_j|\) is large; “All” denotes that we explore all states in \(W_j\), and Rand denotes sampling first 5000 states randomly from \(W_j\) and exploring them. The E, M, and H columns give the number of easy, medium, or hard states among the sampled states. For each \(k = 1, 2, 3\) and the sum of E, M, and H columns is equal to the number of sampled states, and \(\ast\) denotes the number of remaining states.

Observe that \(|V_j|\) is small fraction of \(|V|\) (this illustrates the significance of our use of symbolic methods as opposed to the prohibitive explicit enumeration search). Also, observe that states labeled medium and hard are a small fraction of the sampled states (this illustrates the significance of our efficient iterative sampling strategy).

Table 2: Table 2: CONNECT-3 & -4 against depth-2 strategy of opponent.

<table>
<thead>
<tr>
<th>Game State</th>
<th>j</th>
<th>Win Cond</th>
<th>No. of Sampling</th>
<th>(E_1) = 1</th>
<th>(E_2) = 1</th>
<th>(M_1) = 1</th>
<th>(M_2) = 1</th>
<th>(H_1) = 1</th>
<th>(H_2) = 1</th>
</tr>
</thead>
<tbody>
<tr>
<td>(C-3) 4 \times 4 \times 4 2</td>
<td>RCD</td>
<td>110 All</td>
<td>(24) (9) (0) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(4x4) 6.5 \times 10^2</td>
<td>RC</td>
<td>200 All</td>
<td>(39) (0) (23) (8) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(7.6 \times 10^3)</td>
<td>RD</td>
<td>418 All</td>
<td>(36) (17) (25) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(6.5 \times 10^3)</td>
<td>CD</td>
<td>277 All</td>
<td>(41) (24) (27) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 3: Bottom-2 against depth-3 strategy of opponent.

<table>
<thead>
<tr>
<th>Board State</th>
<th>j</th>
<th>Win Cond</th>
<th>(W_1)</th>
<th>Sampling</th>
<th>(k_1) = 1</th>
<th>(k_2) = 2</th>
<th>(k_3) = 3</th>
</tr>
</thead>
<tbody>
<tr>
<td>3x3 4.1 \times 10^2</td>
<td>2</td>
<td>RCD</td>
<td>20 All</td>
<td>(5) (1) (0) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>RC</td>
<td>0</td>
<td>-</td>
<td>-</td>
<td>-</td>
<td>-</td>
<td>-</td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>RD</td>
<td>9</td>
<td>All</td>
<td>(2) (1) (3) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>CD</td>
<td>1</td>
<td>All</td>
<td>(0) (0) (0) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 4: Bottom-2 against depth-2 strategy of opponent.

<table>
<thead>
<tr>
<th>Board State</th>
<th>j</th>
<th>Win Cond</th>
<th>(W_1)</th>
<th>Sampling</th>
<th>(k_1) = 1</th>
<th>(k_2) = 2</th>
<th>(k_3) = 3</th>
</tr>
</thead>
<tbody>
<tr>
<td>3x3 4.1 \times 10^2</td>
<td>2</td>
<td>RCD</td>
<td>20 All</td>
<td>(5) (0) (1) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>RC</td>
<td>0</td>
<td>-</td>
<td>-</td>
<td>-</td>
<td>-</td>
<td>-</td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>RD</td>
<td>9</td>
<td>All</td>
<td>(2) (1) (3) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4.3 \times 10^3</td>
<td>CD</td>
<td>1</td>
<td>All</td>
<td>(0) (0) (0) (0) (0) (0) (0) (0) (0) (0) (0)</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Description of tables. The caption of Table 1 describes the various column headings used in Tables 1-6, which describe the experimental results. In our experiments, we explore vertices from \(W_2\) and \(W_3\) only as the set \(W_1\) is almost always empty (i.e., if there is a winning starting position it belongs to either \(W_2\) or \(W_3\)). The third column \(j = 2, 3\) denotes whether we explore from \(W_2\) or \(W_3\). For the classification of a given board position as \(E, M, H\), we run the game between the depth-\(k_1\) vs. the depth-\(k_2\) strategy for 30 times. If player 1 (i) wins more than \(\frac{7}{10}\) times (20 times), then the position is identified as easy (E); (ii) wins less than \(\frac{7}{10}\) times (10 times), then it is identified as hard (H); (iii) else as medium (M).

Experimental results for CONNECT games. In Table 1 (resp. Table 2) we present the experimental results for CONNECT-3 and CONNECT-4 games, against depth-3 (resp. depth-2) strategies of the opponent. One of the interesting findings is that in CONNECT-4 games with board size \(5 \times 5\), for all winning conditions (RCD, RD, CD, RC), there are easy, medium, and hard states, for \(k_1 = 1, 2, 3\), when \(j = 3\). In other words, even in much smaller board size \((5 \times 5\) as compared to the traditional \(7 \times 7\)) we discover interesting starting positions for CONNECT-4 games and its simple variants.
Table 5: Tic-Tac-Toe against depth-3 strategy of opponent. The sampling B100 denotes exploring states with the least scored hundred states according to iterative simulation score.

Table 6: Tic-Tac-Toe against depth-2 strategy of opponent.

Experimental results for Bottom-2 games. The results for Bottom-2 (partial gravity-2) against depth-3 (resp. depth-2) strategies of the opponent are shown in Table 3 (resp. Table 4). In contrast to CONNECT games, we observe that medium or hard states do not exist for depth-3 strategies of the player.

Experimental results for Tic-Tac-Toe games. The results for Tic-Tac-Toe games are shown in Table 5 and Table 6. For Tic-Tac-Toe games the strategy exploration is expensive (a tree of depth-3 requires exploration of $10^6$ nodes), and thus exploring many states is time consuming. Hence using the iterative simulation techniques we first assign a score to all states and use exploration for bottom states given certain parameters (namely, expertise levels of players, difficulty level). Existing research on CONNECT games has focussed on establishing which of these games have a winning strategy [9, 21]. In contrast, we show that even simpler variants can be interesting if we start from certain specific states. Existing research has focussed on establishing whether or not these games have a winning strategy.

Summary. We summarize our results in Table 7. We call a state category-i if it is not easy for depth-(i−1) strategy, but it is easy for depth-i strategy. In Table 7 we summarize the different games and the existence of category i states in such games. The generation of $W_j$ for $j = 2$ and $j = 3$ took between two to four hours per game (note that this is a one-time computation for each game). The evaluation time to classify a state as E, M, or H is as follows: for depth-3 strategies of both players, playing 30 times from a board position on average takes (i) 12 seconds for CONNECT-4 games with board size 5 × 5; (ii) 47 seconds for Bottom-2 games with board size 4 × 4; and (iii) 25 minutes for Tic-tac-toe games with board size 4 × 4. Note that the size of the state space is around $10^6$ for CONNECT-4 games with board size 5 × 5, and testing the winning nature of a state takes at least few seconds. Hence an explicit enumeration of the state space cannot be used to obtain $W_j$ in reasonable time; in contrast, our symbolic methods succeed to compute $W_j$ efficiently.

Important findings. We now highlight two important findings of our experimental results.

- Our first key finding is the existence of states of different hardness levels in various games. Let us consider the case where the depth of the opponent strategy is $k_2 = 3$. We observe that in Tic-Tac-Toe games only board positions that are hard for $k_1 = 1$ exist; in particular, and very interestingly, they also exist in board size 4 × 4. Since the default start state in Tic-Tac-Toe games in board size 4 × 4 is heavily biased towards the player who starts first, they have been believed to be uninteresting for ages, whereas our experiments discover interesting starting states for them. With the slight variation of allowable moves (Bottom-2), we obtain board positions that are hard for $k_1 = 2$. In Connect-4 we obtain states that are hard for $k_1 = 3$ even with small board size of 5 × 5.

- The second key finding of our results is the fact that the number of interesting states is a negligible fraction of the huge state space. For example, in Bottom-2 RCD games with board size 4 × 4 the size of the state space is over 1.8 million, but has only two positions that are hard for $k_1 = 2$; and in CONNECT-4 RCD games with board size 5 × 5 the state space is around sixty nine million, but has around two hundred hard states for $k_1 = 3$ and $k_2 = 3$, when $j = 3$, among the five thousand states sampled from $W_j$. Since the size of $W_j$ in this case is around $2.8 \times 10^9$, the total number of hard states is around twelve thousand (among sixty nine million state space size). Since the interesting positions are quite rare, a naive approach of randomly generating positions and measuring its hardness will be searching for a needle in a haystack and be ineffective to generate interesting positions. Thus there is need for a non-trivial search strategy (§3), which our tool implements.

Example board positions. In Figure 1(a)-Figure 1(f) we present examples of several board positions that are of different hardness level for strategies of certain depth. Also see appendix for an illustration. In all the figures, player-X is the current player against opponent of depth-3 strategy. All these board positions were discovered through our experiments.

6. RELATED WORK

Tic-Tac-Toe and Connect-4. Tic-Tac-Toe has been generalized to different board sizes, match length [15], and even polyomino matches [12] to find variants that are interesting from the default start state. Existing research has focussed on establishing whether or not these games have a winning strategy [9, 10, 21]. In contrast, we study how easy or difficult it is to win from winning states given the expertise levels.

Level generation. Our proposed technique, which generates starting states given certain parameters (namely, expertise levels of players, and hardness level), can be used to generate different game levels by simply varying the choice of the parameters along some
partial/total order. Hence, we discuss some related work from the area of level generation.

The problem of level generation has been studied for specific games. Goldspinner [22] is an automatic level generation system for Kgolddrung, which is a puzzle-oriented platform game with dynamic elements. It uses a genetic algorithm to generate candidate levels and simulation to evaluate dynamic aspects of the game. We also use simulation to evaluate the dynamic aspect, but use symbolic methods to generate candidate states; also, our system is parameterized by game rules.

Most other work has been restricted to games without opponent and dynamic content such as Sudoku [13, 23]. Smith et al. used answer set programming to generate levels for the educational puzzle game Refraction that adhered to prespecified constraints written in first-order logic [18]. Similar approaches have also been used to generate levels for platform games [19]. In all these approaches, designers must explicitly specify constraints that the generated content must reflect, for example, the tree needs to be near the rock and the river needs to be near the tree. In contrast, our system takes as input rules of the game and does not require any further help from the designer. [4] also uses a similar model and applies symbolic methods (namely, test input generation techniques) to generate various levels for DragoBox, an algebra-learning video game that became the most purchased game in Norway on the Apple App Store [14]. In contrast, we use symbolic methods for generating valid start states, and use simulation for estimating their hardness level.

**Problem generation.** Recently, there has been interesting work on generating fresh set of practice problems for various subject domains (in the context of intelligent adaptive tutoring) including procedural problems for middle-school mathematics [4], and proof problems for high-school algebra [17] and natural deduction [2]. The work on problem generation for algebraic proof problems [17] uses probabilistic testing to guarantee the validity of a generated problem candidate (from abstraction of the original problem) on random inputs, but there is no guarantee of the hardness level. Our simulation can be likened to this probabilistic testing approach, but it is used to guarantee hardness level; whereas validity is guaranteed by symbolic methods. The work on problem generation for natural deduction [2] involves a backward existential search over the state space of all possible proofs for all possible facts to dish out problems with a specific hardness level. In contrast, we employ a two-phased strategy of backward and forward search; backward search is necessary to identify interesting goals, while forward search ensures hardness levels. Furthermore, our state transitions alternate between different players, thereby necessitating alternate universal vs. existential search over transitions.

Interesting starting states that require few steps to play and win are often published in newspapers and magazines for sophisticated games like Chess and Bridge. These are usually obtained from database of past games. In contrast, we show how to automatically generate such states, albeit for simpler games.

**7. CONCLUSION AND FUTURE WORK**

We revisit the classic domain of simple board games that have been popular since ancient times. We define a novel problem of generating starting states of a given difficulty level, given rules of the game and expertise levels of the players. This offers two key advantages of personalization and freshness. (a) The notion of difficulty is personalized to the players unlike computerized (board) games. There are two parameters that help control the difficulty level: the density of winning game plays and the expertise level of the two players. In contrast, computerized (board) games, do not take into account any notion of expertise level of the player; their notion of difficulty of any level is independent of who is playing the game. (b) It prevents memorization of game plays by introducing freshness in two ways: by presenting multiple start states and by allowing modification of game rules.

We present a novel search technique that combines use of symbolic methods and iterative simulation. For experimental results we focused on traditional board games and their variants. A possible direction of future work would be to investigate methods that can enable our approach to scale to games with more complicated rules. However, as mentioned in §1, finding interesting starting positions in simple and traditional games is both relevant and more challenging. Our experiments identified states of varying difficulty level for various games, opening up new games to be played that were believed to be useless for ages. Moreover as these interesting positions (for traditional games and their variants) are quite rare, their finding is a very useful and important discovery.

8. REFERENCES


Figure 2: Illustration of depth $k_1 = 1$ tree exploration for Tic-Tac-Toe. Since all possible moves for $k_1 = 1$ have equal score of $+0$, all board positions are chosen at random.
Figure 3: Illustration of depth $k_1 = 2$ tree exploration for Tic-tac-toe. In the figure, every choice of player 1 is followed by the optimal choice of opponent and they are collapsed to a single state. Since the center position has a higher score of $+0$, the $k_1 = 2$ strategy always chooses this move and considers all other positions to be equal with score of $-2$. 
Figure 4: Illustration of depth-\(k_1\) strategy exploration on a CONNECT-4 RCD category-3 state. The figure shows how different depth-\(k_1\) strategies choose the best available position to mark on a Connect-4 RCD category-3 state. The example board position of the figure is the same as for Figure 1 (e). The three depth-\(k_1\) strategies (\(k_1 = 1, 2, 3\)) play as player-X and assign a score to each of the three available positions (column-2, 3, 4) by looking \(k_1\)-turns ahead. In each sub-figure, the position with yellow-background is the one chosen for exploration and the positions with grey-background are the predicted moves of how the game might turn out after \(k_1\)-turns. As observed, only \(k_1 = 3\) strategy is able to foresee that marking column-2 would lead player-X to a winning state and also conclude that the other column choices will lead to a draw. Where as, \(k_1 = 1, 2\) incorrectly choose column-3 as the best position to mark hence making this starting position a category-3 state.